

# Digital Number Systems

EE2222 Computer Interfacing and Microprocessors

*Partially based on*

*Microprocessors and Embedded Systems by Hui Wu, UNSW*

*Number Systems by Dr. Paul Beame, University of Washington*

*Number Systems: Negative Numbers by Dr. Chung-Kuan Cheng, UC San Diego*

# Numbers: Positional notation

- Number Base  $b \Rightarrow b$  symbols per digit:
  - Base 10 (Decimal) : 0, 1, 2, 3, 4, 5, 6, 7, 8, 9
  - Base 2 (Binary) : 0, 1
- In general, given a base  $b$ , and number of digits  $p$ , the number is written as

$$d_{p-1}d_{p-2} \dots d_2d_1d_0$$

$$\text{value} = d_{p-1} \times b^{p-1} + d_{p-2} \times b^{p-2} + \dots + d_2 \times b^2 + d_1 \times b + d_0$$

- The leftmost,  $d_{p-1}$ , is the **most significant digit**,  $d_0$ , the rightmost is the **least significant digit**.
  - $1011010 = 1 \times 2^6 + 0 \times 2^5 + 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2 + 0 \times 1 = 64 + 16 + 8 + 2 = 90$

# Digital numbers

- Digital = discrete
  - Binary codes (*example: Binary Coded Decimal*)
- Binary codes
  - Represent symbols using binary digits (bits)
- Digital computers:
  - I/O is digital
    - ASCII, decimal, etc.
  - Internal representation is binary
    - Process information in bits

<u>Decimal Symbols</u>	<u>BCD Code</u>
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

# Digital numbers

- Binary numbers
  - Computers work with 0's and 1's; binary is like the alphabets of a language
- Base conversion
  - For convenience, people use other bases (like octal, decimal, hexadecimal)
  - Need to know how to convert from one to another
- Number systems
  - There are more than one way to express a number in binary.
  - So 1010 could be 10, -2, -5 or -6 and need to know which one.
- A/D and D/A conversion
  - Real world signals come in continuous/analog format
  - It is good to know how they become 0's and 1's (and vice versa)

# Binary Numbers: Base 2

- Bases we will use
  - Binary: Base 2
    - 0,1
  - Octal: Base 8
    - 0,1,2,3,4,5,6,7
  - Decimal: Base 10
    - 0,1,2,3,4,5,6,7,8,9
  - Hexadecimal: Base 16
    - 0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F

- Positional number system
  - $101_2 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$
  - $63_8 = 6 \times 8^1 + 3 \times 8^0$
  - $A1_{16} = 10 \times 16^1 + 1 \times 16^0$
- Addition and subtraction

$$\begin{array}{r} 1011 \\ + 1010 \\ \hline 10101 \end{array}$$

$$\begin{array}{r} 1011 \\ - 0110 \\ \hline 0101 \end{array}$$

# Hexadecimal Numbers: Base 16

- Digits:
  - 0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F
- Normal digits have expected values
- In addition:
  - A → 10
  - B → 11
  - C → 12
  - D → 13
  - E → 14
  - F → 15

# Hexadecimal Numbers: Base 16

- Example (convert hex to decimal):

$$\begin{aligned} \text{B28F0DD} &= (\text{B} \times 16^6) + (2 \times 16^5) + (8 \times 16^4) + (\text{F} \times 16^3) + (0 \times 16^2) + (\text{D} \times 16^1) + (\text{D} \times 16^0) \\ &= (11 \times 16^6) + (2 \times 16^5) + (8 \times 16^4) + (15 \times 16^3) + (0 \times 16^2) + (13 \times 16^1) + (13 \times 16^0) \\ &= 187232477 \text{ decimal} \end{aligned}$$

- Notice that a 7 digit hex number turns out to be a 9 digit decimal number

# Decimal vs. Hexadecimal vs. Binary

- Examples:

- 1010 1100 0101 (binary)  
= ? (hex)

- 10111 (binary)  
= 0001 0111 (binary)  
= ? (hex)

- 3F9(hex)  
= ? (binary)

00	0	0000
01	1	0001
02	2	0010
03	3	0011
04	4	0100
05	5	0101
06	6	0110
07	7	0111
08	8	1000
09	9	1001
10	A	1010
11	B	1011
12	C	1100
13	D	1101
14	E	1110
15	F	1111



# Binary $\rightarrow$ hex/decimal/octal conversion

- Conversion from binary to octal/hex
  - Binary : 10011110001
  - Octal : 10 | 011 | 110 | 001 = 2361<sub>8</sub>
  - Hex : 100 | 1111 | 0001 = 4F1<sub>16</sub>
- Conversion from binary to decimal
  - $101_2 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 = 5_{10}$
  - $63.4_8 = 6 \times 8^1 + 3 \times 8^0 + 4 \times 8^{-1} = 51.5_{10}$
  - $A1_{16} = 10 \times 16^1 + 1 \times 16^0 = 161_{10}$

# Decimal $\rightarrow$ binary/octal/hex conversion

## Binary

	<u>Quotient</u>	<u>Remainder</u>
$56 \div 2 =$	28	0
$28 \div 2 =$	14	0
$14 \div 2 =$	7	0
$7 \div 2 =$	3	1
$3 \div 2 =$	1	1
$1 \div 2 =$	0	1

## Octal

	<u>Quotient</u>	<u>Remainder</u>
$56 \div 8 =$	7	0
$7 \div 8 =$	0	7
$56_{10} = 111000_2$		
$56_{10} = 70_8$		

- Why does this work?
  - $N = 56_{10} = 111000_2$
  - $Q = N/2 = 56/2 = 111000/2 = 11100$  remainder 0
- Each successive divide liberates an LSB (least significant bit)

# Hex → binary conversion

- HEX is a more compact representation of Binary!
- Each hex digit represents 16 decimal values.
- Four binary digits represent 16 decimal values.
- Therefore, each hex digit can replace four binary digits.
- Example:
  - 0011 1011 1001 1010 1100 1010 0000 0000<sub>two</sub>
  - 3 b 9 a c a 0 0<sub>hex</sub>  
C uses notation 0x3b9aca00

# Which Base Should We Use?

- Decimal: Great for humans; most arithmetic is done with these.
- Binary: This is what computers use, so get used to them. Become familiar with how to do basic arithmetic with them (+, -, \*, /).
- Hex: Terrible for arithmetic;
  - *But if we are looking at long strings of binary numbers, it's much easier to convert them to hex in order to look at four bits at a time.*

# How Do We Tell the Difference?

- In general, append a subscript at the end of a number stating the base:
  - $10_{10}$  is in decimal
  - $10_2$  is binary ( $= 2_{10}$ )
  - $10_{16}$  is hex ( $= 16_{10}$ )
- When dealing with AVR microcontrollers:
  - Hex numbers are preceded with “\$” or “0x”
    - $\$10 == 0x10 == 10_{16} == 16_{10}$
  - Binary numbers are preceded with “0b”
  - Octal numbers are preceded with “0” (zero)
  - Everything else by default is Decimal

# Inside the Computer

- To a computer, numbers are always in binary; all that matters is how they are printed out: binary, decimal, hex, etc.
- As a result, it doesn't matter what base a number in C is in...
  - $32_{10} == 0x20 == 100000_2$
- Only the value of the number matters.

# Bits Can Represent Everything

- Characters?
  - 26 letter => 5 bits
  - Upper/lower case + punctuation => 7 bits (in 8) (ASCII)
  - Rest of the world's languages => 16 bits (Unicode)
- Logical values?
  - 0 -> False, 1 => True
- Colors ?
- Locations / addresses? commands?
- But N bits => only  $2^N$  things

# What if too big?

- Numbers really have an infinite number of digits
  - with almost all being zero except for a few of the rightmost digits:
  - e.g: 0000000 ... 000098 == 98
  - Just don't normally show leading zeros
- Computers have fixed number of digits
  - Adding two n-bit numbers may produce an (n+1)-bit result.
  - Since registers' length (8 bits on AVR) is fixed, this is a problem.
  - If the result of add (or any other arithmetic operation), cannot be represented by a register, overflow is said to have occurred



# An Overflow Example

- Example (using 4-bit numbers):

$$\begin{array}{r} +15 \\ \underline{+3} \\ +18 \end{array} \qquad \begin{array}{r} 1111 \\ \underline{0011} \\ 10010 \end{array}$$

- But we don't have room for 5-bit solution, so the solution would be 0010, which is +2, which is wrong.

# How avoid overflow, allow it sometimes?

- Some languages detect overflow (Ada), some don't (C and JAVA)
- Some processors have overflow flags
  - AVR has N, Z, C and V flags to keep track of overflow

# How to Represent Negative Numbers?

- So far, unsigned numbers
- Historically: 3 approaches
  - Sign-and-magnitude
  - Ones-complement
  - Twos-complement
- For all 3, the most-significant bit (MSB) is the sign digit
  - 0  $\equiv$  positive
  - 1  $\equiv$  negative
- Twos-complement is the important one
  - Simplifies arithmetic
  - Used almost universally

# Sign-and-magnitude

- Obvious solution: define leftmost bit to be sign!
- The most-significant bit (MSB) is the sign digit
  - 0  $\equiv$  positive; 1  $\equiv$  negative
- The remaining bits are the number's magnitude
- $+1_{10}$  would be: 0000 0001
- $-1_{10}$  in sign and magnitude would be: 1000 0001

Add		Subtract			Compare and subtract		
4	0100	4	0100	0100	- 4	1100	1100
+ 3	+ 0011	- 3	+ 1011	- 0011	+ 3	+ 0011	- 0011
= 7	= 0111	= 1	$\neq$ 1111	= 0001	- 1	$\neq$ 1111	= 1001

# Shortcomings of sign-and-magnitude

- Problem 1:
  - Two representations for zero
  - $0 = 0000$  and also  $-0 = 1000$
- Problem 2:
  - Arithmetic circuit is more complicated
  - Special steps depending whether signs are the same or not
- Sign and magnitude abandoned because another solution was better

# Ones-complement

- Invert the 0s & 1s of a equivalent binary number provides the 1s complement.
- For positive integer  $x$ , represent  $-x$ :
  - Formula:  $2^n - 1 - x$
  - i.e.  $n=4$ ,  $2^4 - 1 - x = 15 - x$
  - In binary:  $(1\ 1\ 1\ 1) - (b_3\ b_2\ b_1\ b_0)$
  - Just flip all the bits.

# Ones-complement

- Examples:
  - $7_{10} = 00000111_2$
  - $-7_{10} = 11111000_2$
- Questions:
  - What is  $-00000000_2$ ?
  - How many positive numbers in N bits?
  - How many negative numbers in N bits?

# Ones-complement

- Negative number: Bitwise complement positive number
  - $0111 \equiv 7_{10}$
  - $1000 \equiv -7_{10}$

- Solves the arithmetic problem

Add	Invert, add, add carry	Invert and add
$\begin{array}{r} 4 \quad 0100 \\ + 3 \quad + 0011 \\ \hline = 7 \quad = 0111 \end{array}$	$\begin{array}{r} 4 \quad 0100 \\ - 3 \quad + 1100 \\ \hline = 1 \quad 1\ 0000 \\ \text{add carry:} \quad +1 \\ \hline = 0001 \end{array}$	$\begin{array}{r} - 4 \quad 1011 \\ + 3 \quad + 0011 \\ \hline - 1 \quad 1110 \end{array}$

- Remaining problem: Two representations for zero
  - $0 = 0000$  and also  $-0 = 1111$



# Why ones-complement works?

- The ones-complement of an 8-bit positive  $y$  is  $11111111_2 - y$
- What is  $11111111_2$  ?
  - 1 less than  $1\ 00000000_2 \equiv 2^8 \equiv 256_{10}$
  - So in ones-complement  $-y$  is represented by  $(2^8 - 1) - y$
- Adding representations of  $x$  and  $-y$  where  $x, y$  are positive:  
we get  $(2^8 - 1) + x - y$ 
  - If  $x < y$  then  $x - y < 0$  there is no carry and get  $-ve$  number
    - Just add the representations if no carry
  - If  $x > y$  then  $x - y > 0$  there is a carry and get  $+ve$  number
    - Need to add 1 and ignore the  $2^8$ , i.e. “add the carry”
  - If  $x = y$  then answer should be 0, get  $2^8 - 1 = 11111111_2$

# Arithmetic Operations: 1's Complement

Input: two positive integers  $x$  &  $y$ ,

1. We represent the operands in one's complement.
2. We sum up the two operands.
3. We delete  $2^n-1$  if there is carry out at left.
4. The result is the solution in one's complement.

Arithmetic	1's complement
$x + y$	$x + y$
$x - y$	$x + (2^n - 1 - y) = 2^n - 1 + (x - y)$
$-x + y$	$(2^n - 1 - x) + y = 2^n - 1 + (-x + y)$
$-x - y$	$(2^n - 1 - x) + (2^n - 1 - y) = 2^n - 1 + (2^n - 1 - x - y)$

# Arithmetic Operations: Example: $4 - 3 = 1$

$$4_{10} = 0100_2$$

$$3_{10} = 0011_2 \quad -3_{10} \rightarrow 1100_2 \text{ in one's complement}$$

$$\begin{array}{r} 0100 \text{ (4 in decimal )} \\ + 1100 \text{ (12 in decimal or 15-3 )} \\ \hline 1,0000 \text{ (16 in decimal or 15+1 )} \\ 0001 \text{ (after deleting } 2^n - 1) \end{array}$$

We discard the extra 1 at the left which is  $2^n$  and add one at the first bit.

# Arithmetic Operations: Example: $-4 + 3 = -1$

$$4_{10} = 0100_2 \quad -4_{10} \rightarrow \text{Using one's comp.} \rightarrow 1011_2$$

(Invert bits)

$$3_{10} = 0011_2$$

$$\begin{array}{r} 1011 \text{ ( 11 in decimal or 15-4 )} \\ + 0011 \text{ ( 3 in decimal )} \\ \hline 1110 \text{ ( 14 in decimal or 15-1 )} \end{array}$$

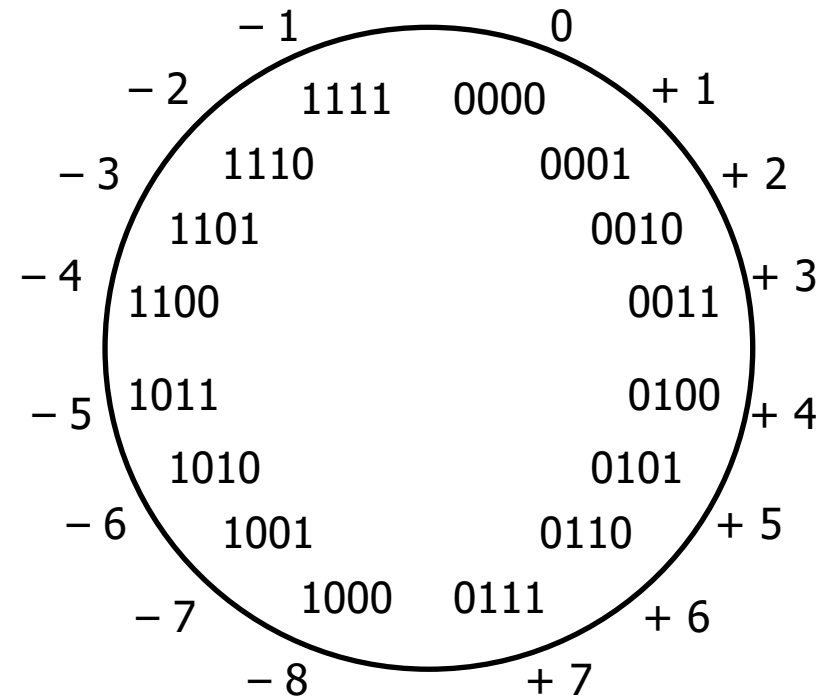
If the left-most bit is 1, it means that we have a negative number.

# Two's complement

- Add 1 to the one's complement provides the two's complement.
- For positive integer  $x$ , represent  $-x$ :
  - Formula:  $2^n - x$
  - i.e.  $n=4$ ,  $2^4 - x = 16 - x$
  - In binary:  $(1\ 0\ 0\ 0\ 0) - (0\ b_3\ b_2\ b_1\ b_0)$
  - Just flip all the bits and add 1.

# Twos-complement

- Negative number: Bitwise complement plus one
  - $0111 \equiv 7_{10}$
  - $1001 \equiv -7_{10}$
- Number wheel
  - Only one zero!
  - MSB is the sign digit
  - $0 \equiv$  positive
  - $1 \equiv$  negative



# Twos-complement

- Complementing a complement  $\square$  the original number
- Arithmetic is easy
  - Subtraction = negation and addition
  - Easy to implement in hardware

Add		Invert and add		Invert and add	
4	0100	4	0100	- 4	1100
+ 3	+ 0011	- 3	+ 1101	+ 3	+ 0011
= 7	= 0111	= 1	1 0001	- 1	1111
		drop carry	= 0001		

# Why twos-complement works better

- Recall:
  - The ones-complement of a b-bit positive  $y$  is  $(2^b-1) - y$
- Adding 1 to get the twos-complement represents  $-y$  by  $2^b - y$ 
  - So  $-y$  and  $2^b - y$  are equal mod  $2^b$   
(leave the same remainder when divided by  $2^b$ )
  - Ignoring carries is equivalent to doing arithmetic mod  $2^b$
- Adding representations of  $x$  and  $-y$  yields  $2^b + x - y$ 
  - If there is a carry then that means  $x \geq y$  and dropping the carry yields  $x-y$
  - If there is no carry then  $x < y$  and then we can think of it as  $2^b - (y-x)$



# Arithmetic Operations: 2's Complement

Input: two positive integers  $x$  &  $y$ ,

1. We represent the operands in two's complement.
2. We sum up the two operands and ignore bit  $n$ .
3. The result is the solution in two's complement.

Arithmetic	2's complement
$x + y$	$x + y$
$x - y$	$x + (2^n - y) = 2^n + (x - y)$
$-x + y$	$(2^n - x) + y = 2^n + (-x + y)$
$-x - y$	$(2^n - x) + (2^n - y) = 2^n + 2^n - x - y$

# Arithmetic Operations: Example: $4 - 3 = 1$

$$4_{10} = 0100_2$$

$$3_{10} = 0011_2 \quad -3_{10} \rightarrow 1101_2$$

$$\begin{array}{r} 0100 \\ + 1101 \\ \hline 10001 \end{array} \rightarrow 1 \text{ (after discarding extra bit)}$$

We discard the extra 1 at the left which is  $2^n$  from 2's complement of -3. Note that bit  $b_{n-1}$  is 0. Thus, the result is positive.

# Arithmetic Operations: Example: $-4 + 3 = -1$

$$4_{10} = 0100_2 \quad -4_{10} \rightarrow \text{Using two's comp.} \rightarrow 1011 + 1 = 1100_2$$

(Invert bits)

$$3_{10} = 0011_2$$

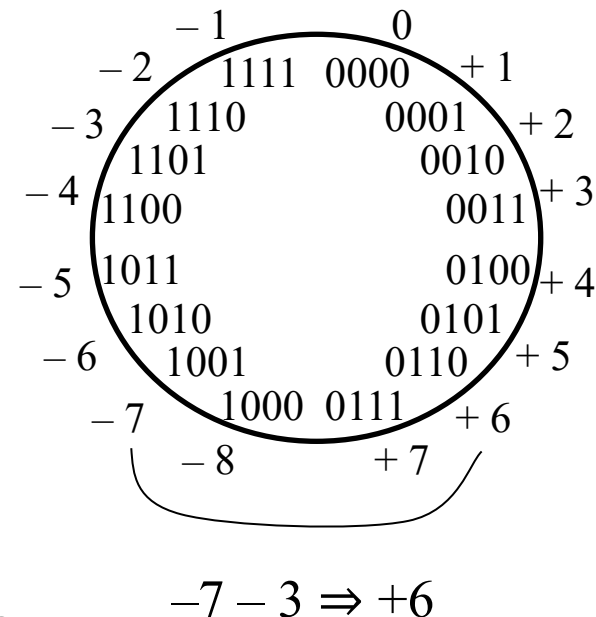
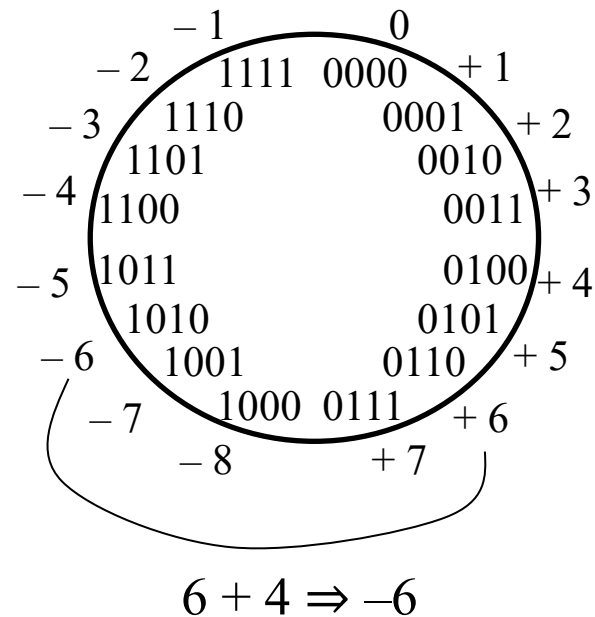
$$\begin{array}{r} 1100 \\ + 0011 \\ \hline \end{array}$$

**1111**  $\rightarrow$  Using two's comp.  $\rightarrow 0000 + 1 = 1$ , so our answer is -1

If left-most bit is 1, it means that we have a negative number.

# Twos-complement overflow

- The rules for detecting overflow in a two's complement sum are simple:
  - Summing two positive numbers can give negative result
  - Summing two negative numbers can give a positive result



# Representing fractional numbers

- To represent fractional numbers, we simply extend the positional system to include digits corresponding to negative powers.
- A number which includes a  $q$  bit fractional part, for a total of  $p+q$  bits:

$$d_{p-1}d_{p-2}\dots d_2d_1d_0.d_{-1}d_{-2}\dots d_{-q+1}d_{-q}$$

- and its value is:

$$value = d_{p-1} \times b^{p-1} + d_{m-2} \times b_{p-2} + \dots + d_2 \times b_2 + d_1 \times b + d_0 + d_{-1} \times b_{-1} + d_{-2} \times b_{-2} + \dots + d_{-q+1} \times b_{-q+1} + d_{-q} \times b_{-q}$$

- Two's complement for fractional numbers:
  - $1.6875_{10} = 01.1011_2$
  - $-1.6875_{10} = 10.0101_2$

# Sign extension

- increasing the number of bits of a binary number while preserving the number's sign (positive/negative) and value
- done by appending digits to the most significant side of the number
- Example:
  - Write +6 and -6 as 2's complement
    - 0110 and 1010
  - Sign extend to 8-bit bytes
    - 00000110 and 11111010

# Still...

- Can't infer a representation from a number
  - 11001 is 25 (unsigned)
  - 11001 is  $-9$  (sign magnitude)
  - 11001 is  $-6$  (ones complement)
  - 11001 is  $-7$  (twos complement)

# Summary

- Positional notation
- Binary/Octal/Decimal/Hexadecimal numbers
- Negative numbers
  - Sign-and-magnitude
  - Ones-complement
  - Twos-complement
- Arithmetic operations
- Representing fractional numbers